

322 Compilers

Why take this course?

- Understanding tools better; what does the compiler really do?
- Computer Engineering & Architecture people: the compiler is your lens to the world
- Phil Greenpun's 10th rule of programming:

Any sufficiently complicated C or Fortran program contains an ad hoc, informally-specified, bug-ridden, slow implementation of half of Common Lisp.

Interpreters vs Compilers

interpreter : program → answer

compiler : program → program
// no answer!

no interpreter ⇒ programs don't run

Why Compile?

Performance. That's the *only* reason.

Why Compile?: Interpreter overhead

```
addl %rax,%rbx
```

vs

```
(define (interp exp)
  (type-case FAE exp
    [num (n) (num n)]
    [add (lhs rhs)
      (let ([lv (interp lhs)]
            [rv (interp rhs)])
        (type-case FAE-Value lv
          [numV (ln)
            (type-case FAE-Value rv
              [numV (rn) (+ ln rn)]
              [else (error 'interp)])])
          [else (error 'interp)]))])
    ...))
```

Why Compile?: Automate Transformations

- Bad maintenance practices, yet profitable transformations

For example unrolling loops; when the chip sees straight-line code it can go faster:

It can “look ahead” and thus make good guesses about what is going to happen next,

filling in caches early, keeping the pipeline full, etc

Why Compile?: Automate Transformations

- Lower-level details are exposed in destination language

For example, variables might live on the stack or in registers; want to use registers as much as possible

Goalposts

Build a compiler accepting a language (L5) that has:

- Higher-order functions
- Safe, mutable arrays
- Arithmetic on (bounded) integers
- Recursive binding form
- Conditionals

and producing x86-64 assembly

Fib in L5

```
(letrec ([fib
          (lambda (n)
            (if (= n 0)
                0
                (if (= n 1)
                    1
                    (+ (fib (- n 1))
                       (fib (- n 2))))))]
  (print (fib 10)))
```

Fib in L4

no more higher-order functions

```
((print (:fib 10))
 (:fib
  (n)
  (if (= n 0)
      0
      (if (= n 1)
          1
          (+ (:fib (- n 1))
              (:fib (- n 2)))))))
```

Fib in L3

every intermediate result has a name

```
((let ([fibten (:fib 10)])
  (print fibten))
 (:fib
 (n)
 (let ([niszero (= n 0)])
  (if niszero
    0
    (let ([nisone (= n 1)])
     (if nisone
       1
       (let ([n1 (- n 1)])
        (let ([fn1 (:fib n1)])
         (let ([n2 (- n 2)])
          (let ([fn2 (:fib n2)])
           (+ fn2
              fn1))))))))))))))
```

Fib in L2

no more nested expressions

```
(:main                                (:fib
0                                     1
0                                     0
(rdi <- 10)                           (cjump rdi = 0 :zero :nonzero)
((mem rsp -8) <- :fr)                  :zero
(call :fib 1)                          (rax <- 0)
:fr                                     (return)
(rdi <- rax)                            :nonzero
(rdi *= 2)                              (cjump rdi = 1 :one :recur)
(rdi += 1)                              :one
(call print 1)                         (rax <- 1)
(return))                              (return)
                                       :recur
                                       (n <- rdi)
                                       (rdi -= 1)
                                       ((mem rsp -8) <- :for)
                                       (call :fib 1)
                                       :for
                                       (result <- rax)
                                       (n -= 2)
                                       (rdi <- n)
                                       ((mem rsp -8) <- :ftr)
                                       (call :fib 1)
                                       :ftr
                                       (rax += result)
                                       (return))
```

Fib in LI

no more variables (just registers)

```
(:main                                (:fib
0                                     1
0                                     2
(rdi <- 10)                           (cjump rdi = 0 :zero :nonzero)
((mem rsp -8) <- :fr)                  :zero
(call :fib 1)                          (rax <- 0)
:fr                                     (return)
(rdi <- rax)                           :nonzero
(rdi *= 2)                             (cjump rdi = 1 :one :recur)
(rdi += 1)                             :one
(call print 1)                         (rax <- 1)
(return))                             (return)
                                       :recur
                                       ((mem rsp 0) <- r12)
                                       ((mem rsp 8) <- r13)
                                       (r12 <- rdi)
                                       (rdi -= 1)
                                       ((mem rsp -8) <- :for)
                                       (call :fib 1)
                                       :for
                                       (r13 <- rax)
                                       (r12 -= 2)
                                       (rdi <- r12)
                                       ((mem rsp -8) <- :ftr)
                                       (call :fib 1)
                                       :ftr
                                       (rax += r13)
                                       (r12 <- (mem rsp 0))
                                       (r13 <- (mem rsp 8))
                                       (return))
```

Implementation/Project overview

L5 → L4 → L3 → L2, each one step

L2 → L1, multiple steps:

- spilling
- graph coloring
- graph construction
- liveness analysis

Speed test

2 assignments per step: tests & implementation

There is no real “late” code

Use any PL you want (learn a new one!)

<http://www.eecs.northwestern.edu/~robby/courses/322-2015-spring/>

Want to pair program? Send me a note (both members),
with the promise on the web page

More admin details, including grading rubric, on website;
read it