

Cost of Substitution

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}))
```

Cost of Substitution

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}})
```

⇒

```
(interp {with {y 2}
           {+ 100 {+ 99 {+ 98 ... {+ y 1}}}}})
```

Cost of Substitution

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 {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}))

⇒

(interp {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y 1}}}}})

⇒

(interp {+ 100 {+ 99 {+ 98 ... {+ 2 1}}}})

Cost of Substitution

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}})
```

⇒

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(interp {with {y 2}
           {+ 100 {+ 99 {+ 98 ... {+ y 1}}}}})
```

⇒

```
(interp {+ 100 {+ 99 {+ 98 ... {+ 2 1}}}})
```

With **n** variables, evaluation will take $O(n^2)$ time!

Deferring Substitution

(interp {with {x 1}
 {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}))



Deferring Substitution

(interp {with {x 1}
 {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}))

⇒

(interp {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}})

x = 1

Deferring Substitution

(interp {with {x 1}
 {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}})

⇒

(interp {with {y 2}
 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}})

⇒

(interp {+ 100 {+ 99 {+ 98 ... {+ y x}}}})

Deferring Substitution

(interp

```
{with {x 1}
  {with {y 2}
    {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}})
```



⇒

(interp

```
{with {y 2}
  {+ 100 {+ 99 {+ 98 ... {+ y x}}}}})
```

x = 1

⇒

(interp

```
{+ 100 {+ 99 {+ 98 ... {+ y x}}}}})
```

y = 2 x = 1

⇒ ... ⇒

(interp

```
y)
```

y = 2 x = 1

Deferring Substitution with the Same Identifier

```
(interp {with {x 1}
           {with {x 2}
                 x} })
```



Deferring Substitution with the Same Identifier

(interp {with {x 1}
 {with {x 2}
 x} })

⇒

(interp {with {x 2}
 x})

Deferring Substitution with the Same Identifier

(interp {with {x 1}
 {with {x 2}
 x} })

⇒

(interp {with {x 2}
 x})

⇒

(interp x)

Deferring Substitution with the Same Identifier

(interp {with {x 1}
 {with {x 2}
 x} })

⇒

(interp {with {x 2}
 x})

⇒

x = 2 x = 1
(interp x)

Always add to start, then always check from start

Deferring Substitution with the Same Identifier

(interp

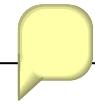
```
{with {x 1}
      (+ {with {x 2}
           x}
          x} )}
```



Deferring Substitution with the Same Identifier

(interp

```
{with {x 1}
      {+ {with {x 2}
            x}
       x} }
```



(interp

```
{+ {with {x 2} x}
    x}
```

x = 1



Deferring Substitution with the Same Identifier

```
(interp {with {x 1}
           (+ {with {x 2}
                  x}
              x} )})
```



```
(interp (+ {with {x 2} x}
            x) )
```

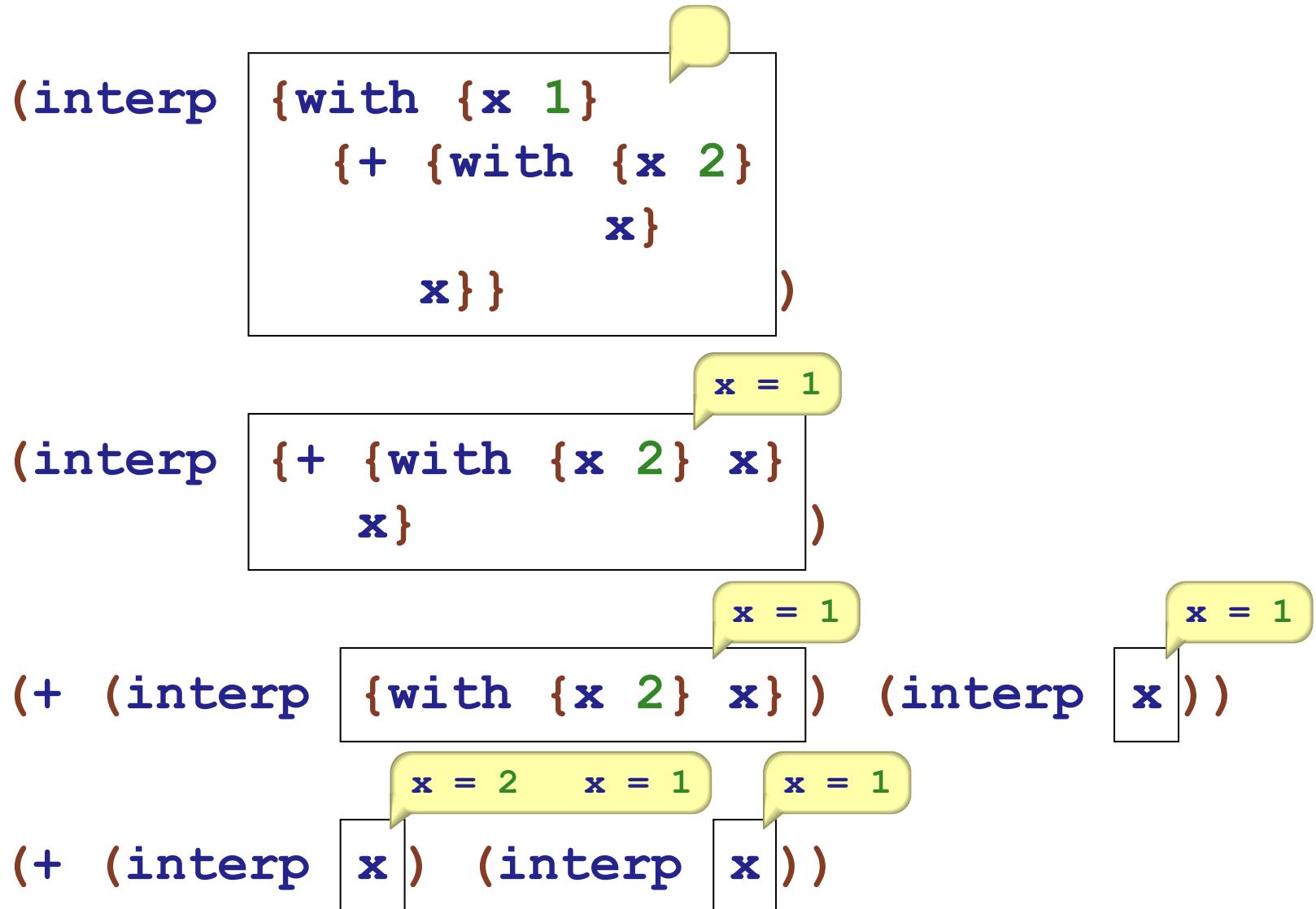


```
(+ (interp {with {x 2} x} ) (interp x) ))
```

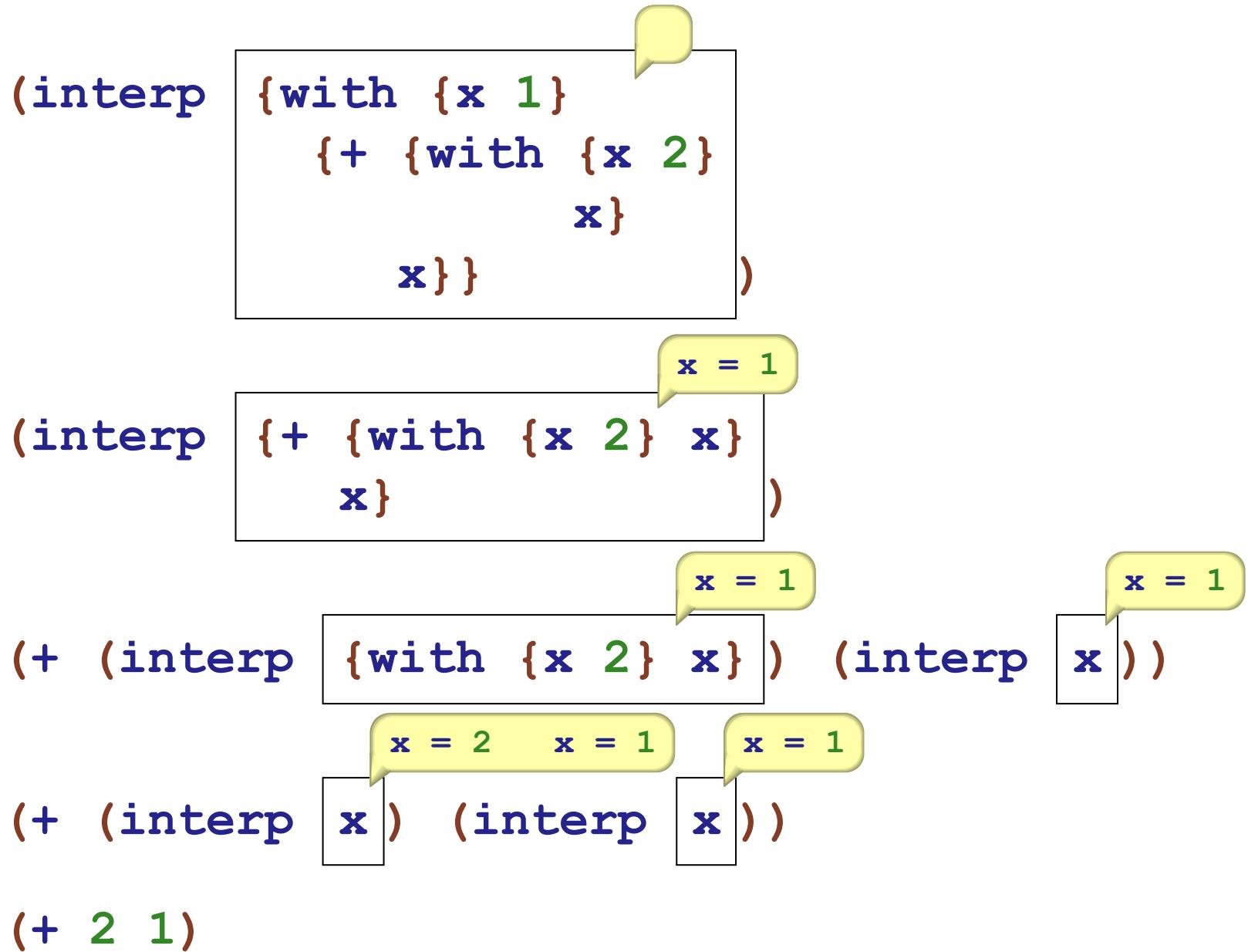
The diagram illustrates the evaluation of a self-referencing lambda expression. It consists of three levels of nested blocks:

- Outermost Level:** Contains the expression `(+ (interp ...))`.
- Middle Level:** Contains the expression `(+ {with {x 2} x})`. A yellow speech bubble labeled `x = 1` is positioned above the `x` in the `{with ...}` block.
- Innermost Level:** Contains the expression `(interp {with {x 1} (+ {with {x 2} x} x)})`. A yellow speech bubble labeled `x = 1` is positioned above the `x` in the `{with {x 1} ...}` block.

Deferring Substitution with the Same Identifier



Deferring Substitution with the Same Identifier



Representing Deferred Substitution

Change

`; interp : WAE -> num`

to

`; interp : WAE DefrdSub -> num`

Representing Deferred Substitution

Change

```
; interp : WAE -> num
```

to

```
; interp : WAE DefrdSub -> num
```

```
(define-type DefrdSub
  [mtSub]
  [aSub (name symbol?)
        (value number?)
        (rest DefrdSub?)] )
```

Interp with DefrdSub

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}}
```

(mtSub))

Interp with DefrdSub

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}}  
(mtSub))
```

```
⇒ (interp {with {y 2}
               {+ 100 {+ 99 {+ 98 ... {+ y x}}}}})  
(aSub 'x 1 (mtSub)))
```

Interp with DefrdSub

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}}  
(mtSub))
```

```
⇒ (interp {with {y 2}
               {+ 100 {+ 99 {+ 98 ... {+ y x}}}}})  
(aSub 'x 1 (mtSub)))
```

```
⇒ (interp {+ 100 {+ 99 {+ 98 ... {+ y x}}}})  
(aSub 'y 2 (aSub 'x 1 (mtSub))))
```

Interp with DefrdSub

```
(interp {with {x 1}
           {with {y 2}
                 {+ 100 {+ 99 {+ 98 ... {+ y x}}}}}}}  
(mtSub))
```

```
⇒ (interp {with {y 2}
           {+ 100 {+ 99 {+ 98 ... {+ y x}}}}})  
(aSub 'x 1 (mtSub)))
```

```
⇒ (interp {+ 100 {+ 99 {+ 98 ... {+ y x}}}})  
(aSub 'y 2 (aSub 'x 1 (mtSub))))
```

⇒ ...

```
⇒ (interp y (aSub 'y 2 (aSub 'x 1 (mtSub))))
```

WAE Interpreter with Deferred Substitutions

```
; interp : WAE DefrdSub -> num
(define (interp a-wae ds)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l ds) (interp r ds))])
    [sub (l r) (- (interp l ds) (interp r ds))])
    [with (bound-id named-expr body-expr)
      ...]
    [id (name) ...]))
```

WAE Interpreter with Deferred Substitutions

```
; interp : WAE DefrdSub -> num
(define (interp a-wae ds)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l ds) (interp r ds))])
    [sub (l r) (- (interp l ds) (interp r ds))])
    [with (bound-id named-expr body-expr)
      ...]
    [id (name) (lookup name ds)])))
```

WAE Interpreter with Deferred Substitutions

```
; lookup : symbol DefrdSub -> num
(define (lookup name ds)
  (type-case DefrdSub ds
    [mtSub () (error 'lookup "free variable")]
    [aSub (sub-name num rest-ds)
      (if (symbol=? sub-name name)
          num
          (lookup name rest-ds))]))
```

WAE Interpreter with Deferred Substitutions

```
; interp : WAE DefrdSub -> num
(define (interp a-wae ds)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l ds) (interp r ds))])
    [sub (l r) (- (interp l ds) (interp r ds))])
    [with (bound-id named-expr body-expr)
      ...]
    [id (name) (lookup name ds)])))
```

WAE Interpreter with Deferred Substitutions

```
; interp : WAE DefrdSub -> num
(define (interp a-wae ds)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l ds) (interp r ds))])
    [sub (l r) (- (interp l ds) (interp r ds))])
    [with (bound-id named-expr body-expr)
          ... (interp named-expr ds) ...])
    [id (name) (lookup name ds)])))
```

WAE Interpreter with Deferred Substitutions

```
; interp : WAE DefrdSub -> num
(define (interp a-wae ds)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l ds) (interp r ds))])
    [sub (l r) (- (interp l ds) (interp r ds))])
    [with (bound-id named-expr body-expr)
      ...
      (aSub bound-id (interp named-expr ds) ds)
      ...]
    [id (name) (lookup name ds)])))
```

WAE Interpreter with Deferred Substitutions

```
; interp : WAE DefrdSub -> num
(define (interp a-wae ds)
  (type-case WAE a-wae
    [num (n) n]
    [add (l r) (+ (interp l ds) (interp r ds))])
    [sub (l r) (- (interp l ds) (interp r ds))])
    [with (bound-id named-expr body-expr)
      (interp
        body-expr
        (aSub bound-id (interp named-expr ds) ds))])
    [id (name) (lookup name ds)])))
```

Function Calls

```
{defun f x) (+ 1 x)}
```

```
(interp (with y 2  
      (f 10)))
```



Function Calls

```
{defun {f x} {+ 1 x}}
```

```
(interp {with {y 2}  
        {f 10}})
```

⇒

```
(interp {f 10})
```

Function Calls

```
{defun {f x} {+ 1 x}}
```

```
(interp {with {y 2}  
        {f 10}})
```

⇒

```
(interp {f 10})
```

y = 2

⇒

```
(interp {+ 1 x})
```

...

Function Calls

```
{defun {f x} {+ 1 x}}
```

```
(interp {with {y 2}  
        {f 10}})
```

⇒

```
(interp {f 10})
```

y = 2

⇒

```
(interp {+ 1 x})
```

x = 10

Interpreting function body starts with only one substitution

Function Calls

What goes wrong if you extend the old substitution?

```
{defun (f x) (+ y x)}
```

```
(interp (with (y 2)
              (f 10)))
```



Function Calls

What goes wrong if you extend the old substitution?

```
{defun {f x} {+ y x}}
```

```
(interp
```

```
{with {y 2}  
{f 10}})
```

⇒

```
(interp
```

```
{f 10})
```

Function Calls

What goes wrong if you extend the old substitution?

```
{defun {f x} {+ y x}}
```

```
(interp
```

```
{with {y 2}  
{f 10}})
```

⇒

```
(interp
```

```
{f 10})
```

⇒

```
(interp
```

```
{+ y x})
```

⇒ 12 wrong!

Function Calls

What goes wrong if you extend the old substitution?

```
{defun {f x} {+ y x}}
```

```
(interp
```

```
{with {y 2}  
{f 10}})
```

⇒

```
(interp
```

```
{f 10})
```

⇒

```
(interp
```

```
{+ y x})
```

⇒ "free var: y"

F1WAE Interpreter with Deferred Substitutions

```
; interp : F1WAE list-of-FunDef DefrdSub -> num
(define (interp a-f1wae fundefs ds)
  (type-case F1WAE a-f1wae
    ...
    [app (name arg-expr)
      ...]))
```

F1WAE Interpreter with Deferred Substitutions

```
; interp : F1WAE list-of-FunDef DefrdSub -> num
(define (interp a-f1wae fundefs ds)
  (type-case F1WAE a-f1wae
    ...
    [app (name arg-expr)
      (local [(define a-fundef
                  (lookup-fundef name fundefs))
              (interp (fundef-body a-fundef)
                     fundefs
                     ...
                     (interp arg-expr fundefs ds)
                     ...))]))
```

F1WAE Interpreter with Deferred Substitutions

```
; interp : F1WAE list-of-FunDef DefrdSub -> num
(define (interp a-f1wae fundefs ds)
  (type-case F1WAE a-f1wae
    ...
    [app (name arg-expr)
      (local [(define a-fundef
                    (lookup-fundef name fundefs))]
        (interp (fundef-body a-fundef)
          fundefs
          (aSub (fundef-arg-name a-fundef)
            (interp arg-expr fundefs ds)
            (mtSub))))]))
```

Timing tests

```
(define (mk-sums n)
  (cond
    [(zero? n) 1]
    [else
      (let ([varn (string->symbol (format "x~a" n))])
        `{+ ,varn ,(mk-sums (- n 1))})))))

(define (mk-withs n body)
  (cond
    [(zero? n) body]
    [else
      (let ([varn (string->symbol (format "x~a" n))])
        `{with {,varn 1}
            ,(mk-withs (- n 1) body)}))))
```

Timing tests, 2

```
(define (mk-exp n) (mk-withs n (mk-sums n)))  
  
(test (mk-exp 2)  
  ` {with {x2 1}  
      {with {x1 1}  
          {+ x2 {+ x1 1} } } } )  
  
(define (run n)  
  (let ([expr (parse (mk-exp n))])  
    (time (interp-expr expr '() ))))
```

Timing tests, 3

With the substitution-based interpreter, expect the difference between adjacent timings to be growing linearly. With the environment-based one, you will also see linear growth, but if you make the environment use a more efficient data structure, that'll go away

(you may need to make the numbers bigger or smaller to see what is going on here)

```
(collect-garbage) (collect-garbage)  
(collect-garbage) (collect-garbage)  
(run 100) (run 110) (run 120)  
(run 130) (run 140) (run 160)
```